

If we approximate the DTFT

$$X(f) = \sum_{n=-\infty}^{\infty} x_n e^{-j2\pi f n}$$

by

$$X(f) \approx \sum_{n=n_1}^{n_2} x_n e^{-j2\pi f n},$$

then the right-hand side can be plotted in MATLAB as follows. Assuming $\mathbf{x} = [x_{n_1}, \dots, x_{n_2}]$,

```
f = linspace(-1/2, 1/2, 201);
nvec = [n1:n2];
x = dtft(f, x, nvec);
plot(f, X)
```

where

```
function Y = dtft(f, x, nvec)
fvec = reshape(f, 1, prod(size(f))); % convert f to row vector
Y = x*exp(-j*2*pi*nvec.*fvec);
Y = reshape(Y, size(f)); % make output have shape of f.
```

will plot an approximation of $X(f)$. Of course, if x_n is a finite-duration signal on $n_1 \leq n \leq n_2$, then `dtft` computes $X(f)$ exactly.

CHAPTER 3

The DFT and the FFT

In Chapter 1, we saw that signal processing for continuous-time, bandlimited waveforms and systems can be accomplished by discrete-time signal processing. However, computers can only evaluate *finite* sums.

Recall that for an infinite sequence x_n , its DTFT is

$$X(f) := \sum_{m=-\infty}^{\infty} x_m e^{-j2\pi f m} \approx \sum_{m=-M_1}^{M_2} x_m e^{-j2\pi f m}$$

for large M_2 and large (negative) M_1 . The sum on the right contains $M_2 - M_1 + 1$ terms. In order to write the finite sum as a sum starting from zero, we can make the change of variable $n = m - M_1$ to get

$$X(f) \approx \sum_{n=0}^{M_2-M_1} x_{n+M_1} e^{-j2\pi f (n+M_1)}.$$

Although the foregoing approximation involves only a finite sum, a computer cannot evaluate it for all values of f in one period, say $|f| \leq 1/2$. Instead, the computer can only evaluate the sum for finitely many values of f . Let $N := M_2 - M_1 + 1$, and let $f = k/N$ to get

$$\begin{aligned} X(k/N) &\approx \sum_{n=0}^{N-1} x_{n+M_1} e^{-j2\pi k(n+M_1)/N} \\ &= e^{-j2\pi k M_1/N} \sum_{n=0}^{N-1} x_{n+M_1} e^{-j2\pi k n/N}. \end{aligned} \quad (3.1)$$

Regarding this last sum as a function of k , observe that it has period N ; i.e., replacing k with $k+N$ does not change the value of the sum. So we only need to evaluate the sum for $k = 0, \dots, N-1$.

3.1. The Discrete Fourier Transform (DFT)

Given a finite sequence y_0, \dots, y_{N-1} , its **discrete Fourier transform (DFT)** is

$$Y_k := \sum_{n=0}^{N-1} y_n e^{-j2\pi k n/N}. \quad (3.2)$$

It is easy to show that Y_k is a periodic function of k with period N . We show below in Section 3.1.7 that the sequence y_n can be recovered from the DFT sequence Y_0, \dots, Y_{N-1} by the **inverse DFT (IDFT)**

$$y_n = \frac{1}{N} \sum_{k=0}^{N-1} Y_k e^{j2\pi k n/N}. \quad (3.3)$$

Of course the right-hand side is a periodic function of n with period N . Hence, although y_n is only defined for $n = 0, \dots, N-1$, we often think of it as being an infinite-duration periodic signal with period N .

3.1.1. Zero Padding

Given a sequence y_0, \dots, y_{M-1} , its DTFT is

$$Y(f) = \sum_{n=0}^{M-1} y_n e^{-j2\pi f n}.$$

Now consider the zero-padded vector

$$z = [y_0, \dots, y_{M-1}, 0, \dots, 0].$$

If the length of z is N , then the DFT of z is

$$\begin{aligned} \sum_{n=0}^{N-1} z_n e^{-j2\pi kn/N} &= \sum_{n=0}^{M-1} y_n e^{-j2\pi kn/N} \\ &= Y(k/N). \end{aligned}$$

In other words, given M data samples y_0, \dots, y_{M-1} , computing a zero-padded DFT gives you samples of $Y(f)$ that are more closely spaced in frequency.

Now suppose that x_n is a causal sequence of infinite duration with DTFT

$$X(f) = \sum_{n=0}^{\infty} x_n e^{-j2\pi f n}.$$

Put

$$X_M(f) := \sum_{n=0}^{M-1} x_n e^{-j2\pi f n}.$$

Then as $M \rightarrow \infty$, $X_M(f) \rightarrow X(f)$. For fixed M , the DFT of $[x_0, \dots, x_{M-1}, 0, \dots, 0]$ zero padded to length N is

$$X_M(k/N) := \sum_{n=0}^{M-1} x_n e^{-j2\pi kn/N} \quad k = 0, \dots, N-1.$$

In other words, for fixed M , if we compute the DFT of $[x_0, \dots, x_{M-1}, 0, \dots, 0]$ with more and more zeros padded, we do *not* get closer to $X(f)$; we get more closely spaced frequency samples of $X_M(f)$. This is illustrated in Figure 3.1.

3.1.2. The Fast Fourier Transform (FFT)

If $y = [y_0, \dots, y_{N-1}]$, then the DFT of y , $Y = [Y_0, \dots, Y_{N-1}]$, can be computed in MATLAB with the command `Y = fft(y)`. Here **FFT** stands for **fast Fourier transform**. The FFT is a special algorithm that computes the DFT very quickly.

Since the DFT is periodic,¹

$$\begin{aligned} Y_{-1} &= Y_{-1+N} = Y_{N-1} \\ Y_{-2} &= Y_{-2+N} = Y_{N-2} \\ &\vdots \\ Y_{-N/2} &= Y_{-N/2+N} = Y_{N/2}. \end{aligned}$$

¹If N is not even, then $N/2$ should be replaced the greatest integer that is less than or equal to $N/2$, which is denoted by $\lfloor N/2 \rfloor$ and is given by `floor` in MATLAB.

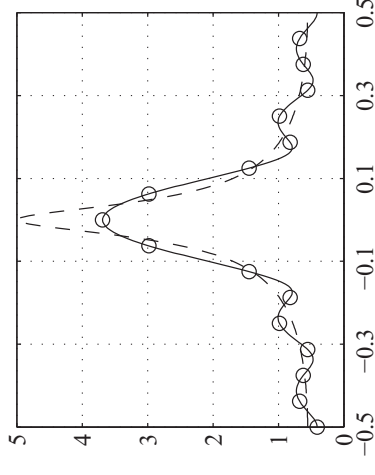


Figure 3.1. The DTFT of an infinite sequence (dashed line), the DFT of the first M elements (solid line), and the DFT of M elements with zero padding (circles).

So, to plot Y_k for $k = -N/2$ to $k = N/2 - 1$, we need to take

$$Y = [Y_0, \dots, Y_{N/2-1}, Y_{N/2}, \dots, Y_{N-1}]$$

$$[Y_{N/2}, \dots, Y_{N-1}, Y_0, Y_1, \dots, Y_{N/2-1}].$$

and convert it to

This is done with the MATLAB command `fftshift`. The corresponding vector of k values can be given by `k = [0 : N-1] - N/2`, and we could then use the command `plot(k, fftshift(Y))`.

3.1.3. Using the FFT to Compute the DFT

If we are approximating the sum in (3.1), then we would use k/N to have the horizontal axis run from $-1/2$ to $1/2$. The MATLAB function `dtftfft` given below computes the right-hand side of (3.1) using `fft` and also takes care of the bookkeeping to return to you the corresponding frequencies f in $[-1/2, 1/2]$. The command for this is `[Y, F] = dtftfft(x, M1)`, where the elements of x are $[x_{M_1}, \dots, x_{M_2}]$. If M_1 is zero, it can be omitted and you can write `[Y, F] = dtftfft(x)` instead. There is also an optional third argument if you want to force `fft` to zero-pad x . Remember, the spacing between frequencies in the DFT and the FFT is $1/\text{length}(x)$. The required command is `[Y, F] = dtftfft(x, M1, N)` to force `dtftfft` to add enough zeros to x to make its length N . If x is longer than N , x will be truncated to N elements. Finally, a nonpositive value of N causes `dtftfft` to zero-pad x so its length is a power of 2. This makes `fft` as fast as possible. Here is the function.

```
function [y,f] = dtftfft(x,varargin)
%dtftfft Compute the discrete-time Fourier transform using FFT.
N = length(x);
if nargin==3
    N = varargin{2};
    if N<=0 % If N<=0, zero pad x to a power of 2.
        N = 2^ceil(log2(length(x)));
    end
end
Y = fft(x,N);
k = [0:N-1];
if nargin>=2
    M1 = varargin{1};
    if M1~=0
        Y = exp(-j*2*pi*M1/N*k) .* Y;
    end
end
Y = fftshift(Y);
% For N=2m even, change 0,...,m,...,2m-1
% to -m,...,0,...,m-1. For N=2m+1 odd, change
% 0,...,m-1,m,m+1,...,2m to -m,...,-1,0,1,...,m.
% Then divide by N to get frequencies in [-1/2,1/2].
f = (k-floor(N/2))/N;
```